Optimization for data science Constrained Optimization and Standard Problems

R. Flamary

Master Data Science, Institut Polytechnique de Paris

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Full course overview

1. Introduction to optimization for data science

- $1.1\,$ ML optimization problems and linear algebra recap
- 1.2 Optimization problems and their properties (Convexity, smoothness)

2. Smooth optimization : Gradient descent

2.1 First order algorithms, convergence for smooth and strongly convex functions

3. Smooth Optimization : Quadratic problems

- 3.1 Solvers for quadratic problems, conjugate gradient
- 3.2 Linesearch methods

4. Non-smooth Optimization : Proximal methods

- 4.1 Proximal operator and proximal algorithms
- 4.2 Lab 1: Lasso and group Lasso

5. Stochastic Gradient Descent

- 5.1 SGD and variance reduction techniques
- 5.2 Lab 2: SGD for Logistic regression

6. Constrained Optimization and Standard Problems

6.1 Lagrangian, LP, QP and Mixed Integer Programming

7. Coordinate descent

7.1 Algorithms and Labs

8. Newton and quasi-newton methods

8.1 Second order methods and Labs

9. Beyond convex optimization

9.1 Nonconvex reg., Frank-Wolfe, DC programming, autodiff

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2. Smooth optimization : Gradient descent		
3. Smooth Optimization : Quadratic problems		
4. Non-smooth optimization : Proximal methods		
5. Stochastic Gradient Descent		
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Constrained optimization and standard problems

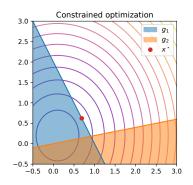
Optimization problem (standard form)

$\min_{\mathbf{x}\in\mathbb{R}^n}$	$F(\mathbf{x})$		
with	$h_j(\mathbf{x}) = 0$	$\forall j = 1, \dots, p$	(1)
and	$g_i(\mathbf{x}) \le 0$	$\forall i=1,\ldots,q.$	

- F is convex and differentiable
- *h_j* and *g_i* are differentiable and define convex constraints.
- The problem is convex.

Standard problems for specialized solvers

- Linear Programming (LP)
- Quadratic Programming (QP)
- Mixed Integer Programming (MIP)
- Quadratically Constrained QP, Second Order Cone Programming, etc.



Constrained optimization and standard problems

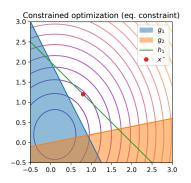
Optimization problem (standard form)

$$\begin{array}{ll} \min_{\mathbf{x}\in\mathbb{R}^n} & F(\mathbf{x}) \\ \text{with} & h_j(\mathbf{x}) = 0 \quad \forall j = 1,\dots,p \quad (1) \\ \text{and} & g_i(\mathbf{x}) \le 0 \quad \forall i = 1,\dots,q. \end{array}$$

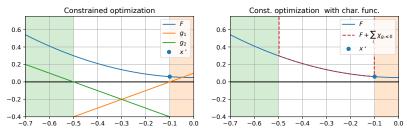
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Standard problems for specialized solvers

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- Quadratic Programming (QP)
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- Quadratically Constrained QP, Second Order Cone Programming, etc.



From indicator to Lagrangian (1)



Constrained optimization problem

$$\min_{\mathbf{x}\in\mathbb{R}^n} F(\mathbf{x})
\text{with} \quad q_i(\mathbf{x}) < 0 \quad \forall i = 1, \dots, q.$$
(2)

The problem is equivalent to

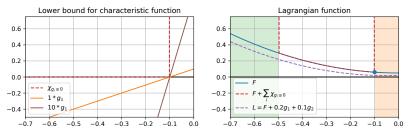
$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad F(\mathbf{x}) + \sum_i \chi_{\leq 0}(g_i(\mathbf{x}))$$

where $\chi_{\leq 0}$ is the characteristic function of the negative values (+ ∞ for positive value and 0 everywhere else).

The problem is not differentiable but can sometimes be solved using :

- Projected Gradient Descent (PGD) for simple constraints.
- Proximal splitting methods for more complex constraints (needs prox.op.).

From indicator to Lagrangian (2)



Lagrangian (details in [Knowles, 2010])

$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad F(\mathbf{x}) + \sum_i \chi_{\leq 0}(g_i(\mathbf{x}))$$

▶ Let us define the Lagrangian function $\mathcal{L}(\mathbf{x}, \mathbf{u})$ such that $\mathbf{u} \ge 0$ and

$$\mathcal{L}(\mathbf{x}, \mathbf{u}) = F(\mathbf{x}) + \sum_{i=1}^{k} u_i g_i(\mathbf{x})$$

For $\mathbf{u} \ge 0$ we have

$$F(\mathbf{x}) + \sum_{i} \chi_{\leq 0}(g_i(\mathbf{x})) = \max_{\mathbf{u} \geq 0} \mathcal{L}(\mathbf{x}, \mathbf{u}) \geq \mathcal{L}(\mathbf{x}, \mathbf{u})$$

Because $\max_{u\geq 0} uv = \chi_{\leq 0}(v)$ that is is $+\infty$ for v > 0 and 0 for $v \leq 0$.

6.1.1 - Constrained optimization - Lagrangian of a constrained optimization problem - 6/52

Lagrangian for standard constrained optimization

Optimization problem (standard form)

$$\min_{\mathbf{x} \in \mathbb{R}^n} \quad F(\mathbf{x})$$
with $h_j(\mathbf{x}) = 0 \quad \forall j = 1, \dots, p$
and $g_i(\mathbf{x}) \le 0 \quad \forall i = 1, \dots, q.$

$$(3)$$

Lagrangian of the optimization problem

We define the Lagrangian of the problem the function $\mathcal L$ such that :

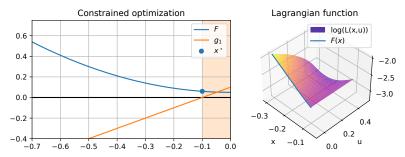
$$\mathcal{L}(\mathbf{x}, \mathbf{u}, \mathbf{v}) = F(\mathbf{x}) + \sum_{i=1}^{k} u_i g_i(\mathbf{x}) + \sum_{j=1}^{m} v_j h_j(\mathbf{x})$$
(4)

- lnequality constraints are multiplied by positive dual variables $u_i \ge 0$.
- Equality constraints are multiplied by dual variables v_j (any sign).
- The constrained optimization problem can be reformulated as

$$\min_{\mathbf{x}} \max_{\mathbf{u} \ge 0, \mathbf{v}} \mathcal{L}(\mathbf{x}, \mathbf{u}, \mathbf{v})$$

Finding a saddle point of the problem solves the constrained problem.

Example of Lagrangian



Optimization problem

$$\min_{x \le -0.1} \quad F(x) = x^2 + 0.05$$

• Constraint
$$g(x) = x + 0.1 \le 0.$$

The Lagrangian is

$$\mathcal{L}(x, u) = x^2 + 0.05 + u(x + 0.1)$$

with u > 0.

6.1.1 - Constrained optimization - Lagrangian of a constrained optimization problem - 8/52

Lagrange dual function

Lagrange dual function

The Lagrange dual function D of the problem is

$$D(\mathbf{u}, \mathbf{v}) = \inf_{\mathbf{x}} \mathcal{L}(\mathbf{x}, \mathbf{u}, \mathbf{v})$$
(5)

If F is not bounded below,
$$D = -\infty$$
.

D is always concave (even when F is non-convex)

Lower bound

For all $\mathbf{u} \geq \mathbf{0}, \mathbf{v}$ and feasible \mathbf{x} we have

$$F(\mathbf{x}) \ge \mathcal{L}(\mathbf{x}, \mathbf{u}, \mathbf{v}) \ge D(\mathbf{u}, \mathbf{v})$$

Proof:

$$\mathcal{L}(\mathbf{x}, \mathbf{u}, \mathbf{v}) = F(\mathbf{x}) + \sum_{i=1}^{k} \underbrace{u_i g_i(\mathbf{x})}_{\leq 0} + \sum_{j=1}^{m} \underbrace{v_j h_j(\mathbf{x})}_{=0} \leq F(\mathbf{x})$$

because x feasible $(g_i(\mathbf{x}) \leq 0, h_j(\mathbf{x}) = 0)$ and $\mathbf{u} \geq 0$.

$$\min_{x,x \ge 0} \quad F(x) = (x+1)^2$$

1. Express the Lagrangian of the problem above :

 $\mathcal{L}(x,u) =$

2. Solve the infimum *w.r.t. x* for a given dual variable *u*:

 $x^{\star} =$

3. Express the Lagrange Dual function D(u):

$$D(u) = \mathcal{L}(x^{\star}, u) =$$

$$\min_{x,x \ge 0} \quad F(x) = (x+1)^2$$

1. Express the Lagrangian of the problem above :

$$\mathcal{L}(x,u) = (x+1)^2 - ux$$

2. Solve the infimum *w.r.t. x* for a given dual variable *u*:

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1. Express the Lagrangian of the problem above :

$$\mathcal{L}(x,u) = (x+1)^2 - ux$$

2. Solve the infimum *w.r.t. x* for a given dual variable *u*:

$$x^{\star} = \frac{u}{2} - 1$$

3. Express the Lagrange Dual function D(u):

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$$\min_{x,x \ge 0} \quad F(x) = (x+1)^2$$

1. Express the Lagrangian of the problem above :

$$\mathcal{L}(x,u) = (x+1)^2 - ux$$

2. Solve the infimum *w.r.t.* x for a given dual variable u:

$$x^{\star} = \frac{u}{2} - 1$$

3. Express the Lagrange Dual function D(u):

$$D(u) = \mathcal{L}(x^*, u) = -\frac{u^2}{4} - u + 4$$

Duality Gap and Strong duality

Definition

For a feasible primal variable ${\bf x}$ and feasible dual variables ${\bf u}, {\bf v}$ we call duality gap the following positive value

$$F(\mathbf{x}) - D(\mathbf{u}, \mathbf{v}) \ge 0 \tag{6}$$

One property of the duality gap is that

$$F(\mathbf{x}) - F^{\star} \le F(\mathbf{x}) - D(\mathbf{u}, \mathbf{v})$$

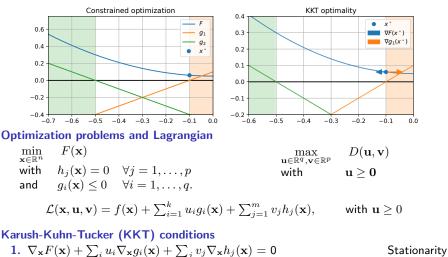
- If the duality gap is 0 for a feasible triplet x^{*}, u^{*}, v^{*} then x^{*} is optimal for the primal and u^{*}, v^{*} are optimal for the dual problem.
- If F^{*} = D^{*} the problem is said to have strong duality.
- Slater's constraint qualification: if the primal problem is convex and there exists a feasible solution :

$$\exists \mathbf{x} \in \mathbb{R}^n, \ h_j(\mathbf{x}) = 0, \ g_i(\mathbf{x}) < 0 \ \forall i, j$$

then strong duality holds.

Karush-Kuhn-Tucker (KKT) conditions

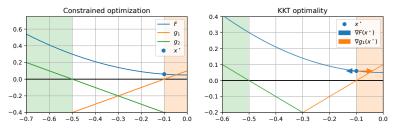
4. $u_i q_i(\mathbf{x}) = 0 \ \forall i$



2. $g_i(\mathbf{x}) \le 0, h_j(\mathbf{x}) = 0, \ \forall i, \forall j$ Primal feasibility**3.** $u_i > 0 \ \forall i$ Dual feasibility

Complementarity

Solution and optimality conditions



Solution of the optimization problem

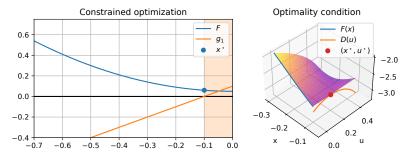
For a problem with strong duality (satisfying Slater's conditions) the two following statements are equivalent:

- $\blacktriangleright~{\bf x}^{\star}$ and ${\bf u}^{\star}, {\bf v}^{\star}$ are solutions of the primal and dual problems.
- $\blacktriangleright \ \mathbf{x}^{\star}$ and $\mathbf{u}^{\star}, \mathbf{v}^{\star}$ satisfy the KKT conditions.

Finding a solution (sometimes)

- 1. Express the Lagrangian and the KKT conditions.
- 2. Try to find an analytic solution for \mathbf{x}^{\star} as function of $\mathbf{u},\mathbf{v}.$
- 3. Express the dual problem and solve it if easier than primal.
- 4. Use KKT to recover the primal solution \mathbf{x}^{\star}

Example of KKT optimality conditions



Optimization problem

$$\min_{x \le -0.1} \quad F(x) = x^2 + 0.05$$

• The Lagrangian is $\mathcal{L}(x, u) = x^2 + 0.05 + u(x + 0.1)$ with $u \ge 0$.

•
$$\nabla_x \mathcal{L}(x, u) = 2x + u = 0$$
 gives $x^* = -\frac{u}{2}$

•
$$D(u) = \mathcal{L}(x^*, u) = -\frac{u^2}{4} + 0.1u + 0.05$$

• Optimal values are $(x^{\star}, u^{\star}) = (-0.1, 0.2).$

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{x}\|^2 \qquad \text{subject to } \sum_{i=1}^n x_i = 1$$

1. Express the Lagrangian of the problem above :

 $\mathcal{L}(x,v) =$

2. Express the KKT of the problem:

2.1
 2.2
 2.3
 2.4

3. Deduce from 1 and 2 above the optimal v^* by maximizing D(v) then \mathbf{x}^* :

$$D(v) =$$

$$v_i^\star = x_i^\star =$$

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{x}\|^2 \qquad \text{subject to } \sum_{i=1}^n x_i = 1$$

1. Express the Lagrangian of the problem above :

$$\mathcal{L}(x,v) = \frac{1}{2} \|\mathbf{x}\|^2 + v(\sum_{i=1}^n x_i - 1)$$

- 2. Express the KKT of the problem:
 - 2.1
 - 2.2
 - 2.3
 - 2
 - 2.4
- **3.** Deduce from 1 and 2 above the optimal v^* by maximizing D(v) then \mathbf{x}^* :

$$D(v) =$$

$$v_i^\star = x_i^\star =$$

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{x}\|^2 \qquad \text{subject to } \sum_{i=1}^n x_i = 1$$

1. Express the Lagrangian of the problem above :

$$\mathcal{L}(x,v) = \frac{1}{2} \|\mathbf{x}\|^2 + v(\sum_{i=1}^n x_i - 1)$$

- 2. Express the KKT of the problem:
 - **2.1** $x_i + v = 0$, $\forall i$ **2.2** $\sum_{i=1}^{n} x_i - 1 = 0$ **2.3** None **2.4** None

3. Deduce from 1 and 2 above the optimal v^* by maximizing D(v) then \mathbf{x}^* :

$$D(v) =$$

$$v_i^\star = x_i^\star =$$

$$\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{x}\|^2 \qquad \text{subject to } \sum_{i=1}^n x_i = 1$$

1. Express the Lagrangian of the problem above :

$$\mathcal{L}(x,v) = \frac{1}{2} \|\mathbf{x}\|^2 + v(\sum_{i=1}^n x_i - 1)$$

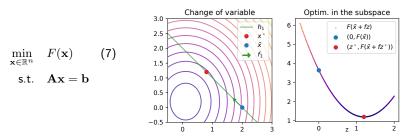
- 2. Express the KKT of the problem:
 - **2.1** $x_i + v = 0$, $\forall i$ **2.2** $\sum_{i=1}^{n} x_i - 1 = 0$ **2.3** None **2.4** None

3. Deduce from 1 and 2 above the optimal v^* by maximizing D(v) then \mathbf{x}^* :

$$D(v) = -\frac{nv^2}{2} - v$$

 $v_i^{\star} = -\frac{1}{n}, \forall i$ $x_i^{\star} = \frac{1}{n}, \forall i$

Linear equality constraints (change of variables)



• With $\mathbf{A} \in \mathbb{R}^{p \times n}$ defining p linearly independent constraints.

▶ We can eliminate the equality constraints using basic linear algebra.

$$\{\mathbf{x}|\mathbf{A}\mathbf{x}=\mathbf{b}\} = \{\mathbf{F}\mathbf{z}+\hat{\mathbf{x}}|\mathbf{z}\in\mathbb{R}^{n-p}\}$$

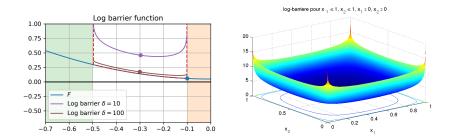
where $\hat{\mathbf{x}}$ is a vector satisfying $\mathbf{A}\hat{\mathbf{x}} = \mathbf{b}$ and $Im(\mathbf{F}) = Ker(\mathbf{A})$.

- ▶ In Python one can compute F with scipy.linalg.null_space.
- The equivalent unconstrained problem is then

$$\min_{\mathbf{z}\in\mathbb{R}^{n-p}} F(\mathbf{F}\mathbf{z}+\hat{\mathbf{x}}) \tag{8}$$

where we can recover the solution of (7) with $\mathbf{x}^{\star}=\mathbf{F}\mathbf{z}^{\star}+\hat{\mathbf{x}}.$

Log-Barrier function



Approximating the inequality constraints

- The **log-barrier** function is an approximation of the characteristic function χ .
- The hard constraints can then be replaced by the log-barrier with $\delta > 0$

$$\begin{array}{ll} \min_{\mathbf{x} \in \mathbb{R}^n} & F(x) \\ \text{s.t.} & g_i(\mathbf{x}) \le 0 \; \forall i \end{array} \qquad \Rightarrow \qquad \min_{\mathbf{x} \in \mathbb{R}^n} & F(\mathbf{x}) + \frac{1}{\delta} \sum_{i=1}^q -\log(-g_i(\mathbf{x})) \end{array}$$

Interior point solver

$$\mathbf{x}(\delta) = \arg\min_{\mathbf{x}\in\mathbb{R}^n} \quad F(\mathbf{x}) + \frac{1}{\delta} \sum_{i=1}^q -\log(-g_i(\mathbf{x})) \tag{9}$$

Interior Point algorithm

Initialize with a feasible ${\bf x},$ and $\delta>0, \mu>1$

1.
$$\mathbf{x} = \mathbf{x}(\delta)$$

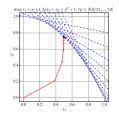
2.
$$\delta = \mu \delta$$

3. Go to 1. until convergence.

Properties of the algorithm

- Requires a solver for the inner problem : computing $\mathbf{x}(\delta)$
- Inner problem is unconstrained and smooth inside the domain.
- Converges to the solution of the constrained problem : $\lim_{\delta \to \infty} \mathbf{x}(\delta) = \mathbf{x}^{\star}$
- All iterations are inside the constraints.
- Converges provably in polynomial time for LP and QP.

More details: [Boyd and Vandenberghe, 2004, Ch.11], [Nocedal and Wright, 2006, Ch. 19] 6.1.3 - Constrained optimization - Interior point solvers - 18/52



Linear Program (LP)

Linear program in standard form

$$\min_{\mathbf{x}\in\mathbb{R}^n} \mathbf{c}^\top \mathbf{x} \tag{10}$$

s.t.
$$Ax = b$$

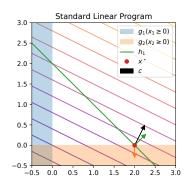
 $\mathbf{x} \ge 0$

- Linear objective with $\mathbf{c} \in \mathbb{R}^n$
- Linear equality constraints with $\mathbf{A} \in \mathbb{R}^{p \times n}, \mathbf{b} \in \mathbb{R}^{p}$
- Positivity inequality constraints.

Problem as a function of $\mathbf{A}\mathbf{x} = \mathbf{b}$

- Underdetermined (p < d): more variables than equations.
- Determined (p = d) : as many independent equations than variables, a unique solution x^{*} = A⁻¹b if A invertible.
- Overdetermined (p > d) : not feasible.

We look at the case where p < d.



Linear Program (LP)

General formulation for LP

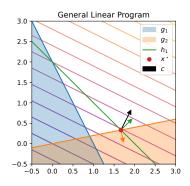
$$\min_{\mathbf{x} \in \mathbb{R}^n} \mathbf{c}^T \mathbf{x}$$
(11)
s.t. $\mathbf{G} \mathbf{x} \le \mathbf{h}$
 $\mathbf{A} \mathbf{x} = \mathbf{b}$

- Closer formulation to the constrained optimization (1).
- $\mathbf{A} \in \mathbb{R}^{p \times n}, \mathbf{b} \in \mathbb{R}^{p}$, and $\mathbf{G} \in \mathbb{R}^{q \times n}, \mathbf{h} \in \mathbb{R}^{q}$.
- Most standard solvers (open source and commercial) use this formulation.

Exercise 3: Classical constraints

Express the matrices and vectors from general LP above for the following constraints:

- Positivity $\mathbf{x} \ge \mathbf{0}$:
- Simplex $\{\mathbf{x} | \mathbf{x} \ge \mathbf{0}, \sum_i x_i = 1\}$:
- Box constraints $l \le x \le u$:



Linear Program (LP)

General formulation for LP

$$\min_{\mathbf{x}\in\mathbb{R}^n} \mathbf{c}^T \mathbf{x}$$
(11)

s.t.
$$\mathbf{G}\mathbf{x} \leq \mathbf{h}$$

 $\mathbf{A}\mathbf{x} = \mathbf{b}$

•
$$\mathbf{A} \in \mathbb{R}^{p \times n}, \mathbf{b} \in \mathbb{R}^{p}$$
, and
 $\mathbf{G} \in \mathbb{R}^{q \times n}, \mathbf{h} \in \mathbb{R}^{q}$.

 Most standard solvers (open source and commercial) use this formulation.

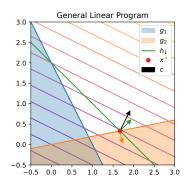
Exercise 3: Classical constraints

Express the matrices and vectors from general LP above for the following constraints:

• Positivity $\mathbf{x} \ge \mathbf{0}$: $\mathbf{G} = -\mathbf{I}_n, \mathbf{h} = \mathbf{0}$

▶ Simplex
$$\{ \mathbf{x} | \mathbf{x} \ge \mathbf{0}, \sum_{i} x_{i} = 1 \}$$
 : $\mathbf{A} = [1, ..., 1], \mathbf{b} = [1], \mathbf{G} = -\mathbf{I}_{n}, \mathbf{h} = \mathbf{0}$

• Box constraints $\mathbf{l} \leq \mathbf{x} \leq \mathbf{u}$: $\mathbf{G} = \begin{bmatrix} \mathbf{I}_n \\ -\mathbf{I}_n \end{bmatrix}$, $\mathbf{h} = \begin{bmatrix} \mathbf{u} \\ \mathbf{l} \end{bmatrix}$



6.2.1 - Linear Program (LP) - Problem formulation - 20/52

Example of LP : Optimal Transport (OT)

Definition of the problem

- *n* factories produce $a_i, \forall i$ amount of goods (vector **a**).
- d stores need to sell $s_j, \forall j$ amount of goods ((vector s, same total as a)).
- There is a cost C_{i,j} of transporting a unitary amount of good from factory i to store j.
- Find the optimal (cheapest) way to move all the goods between factories and stores. A solution of the problem is called a transport matrix.

Optimal transport problem

$$\begin{split} \min_{\mathbf{X} \in \mathbb{R}^{n \times d}} & \sum_{i=1,j=1}^{n,d} C_{i,j} X_{i,j} \\ \text{s.t.} & \sum_{j} X_{i,j} = a_i \quad \forall i, \sum_{i} X_{i,j} = s_j \quad \forall j \\ & X_{i,j} \geq 0 \qquad \forall i,j \end{split}$$

- Resource allocation problem .
- Proposed by [Kantorovich, 1942].
- Nobel prize in economy.
- Now used a lot in machine learning.

Exercise 4: OT expressed as general LP problem

We express the matrix \mathbf{x} as the concatenation of the rows of the matrix \mathbf{X} :

$$\mathbf{x} = [X_{1,1}, X_{1,2}, X_{1,3}, \dots X_{n,d-1}, X_{n,d}]^T$$

The cost matrix \mathbf{C} is also vectorized as \mathbf{c} .

1. Express the row-wise equality constraints $\sum_{i} X_{i,j} = a_i, \forall i \text{ and } \mathbf{A}_1 \mathbf{x} = \mathbf{a}$:

$$\mathbf{A}_1 =$$

The matrix can be expressed simply with the Kroenecker product \otimes

2. Express the column-wise equality constraints $\sum_{i} X_{i,j} = s_j, \forall j \text{ and } \mathbf{A}_2 \mathbf{x} = \mathbf{s}$:

$$\mathbf{A}_2 =$$

3. Express all the matrices in the general LP :

$$\mathbf{A} = \begin{bmatrix} \mathbf{A}_1 \\ \mathbf{A}_2 \end{bmatrix}, \quad \mathbf{b} = , \quad \mathbf{G} = , \quad \mathbf{h} =$$

Exercise 4: OT expressed as general LP problem

We express the matrix ${\bf x}$ as the concatenation of the rows of the matrix ${\bf X}:$

$$\mathbf{x} = [X_{1,1}, X_{1,2}, X_{1,3}, \dots X_{n,d-1}, X_{n,d}]^T$$

The cost matrix \mathbf{C} is also vectorized as \mathbf{c} .

1. Express the row-wise equality constraints $\sum_{i} X_{i,j} = a_i, \forall i \text{ and } \mathbf{A}_1 \mathbf{x} = \mathbf{a}$:

$$\mathbf{A}_1 = \mathbf{I}_n \otimes \mathbf{1}_{1,d}$$

The matrix can be expressed simply with the Kroenecker product \otimes

2. Express the column-wise equality constraints $\sum_{i} X_{i,j} = s_j, \forall j \text{ and } \mathbf{A}_2 \mathbf{x} = \mathbf{s}$:

$$\mathbf{A}_2 = \mathbf{1}_{1,n} \otimes \mathbf{I}_d$$

3. Express all the matrices in the general LP :

$$\mathbf{A} = \begin{bmatrix} \mathbf{A}_1 \\ \mathbf{A}_2 \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} \mathbf{a} \\ \mathbf{s} \end{bmatrix}, \quad \mathbf{G} = -\mathbf{I}_{nd}, \quad \mathbf{h} = \mathbf{0}_{nd}$$

Reduction from general to standard problem

Reformulation to standard LP with positive variables

$$\min_{\mathbf{x}\in\mathbb{R}^n} \mathbf{c}^T \mathbf{x} = \tilde{\mathbf{c}}^\top \tilde{\mathbf{x}}$$
s.t. $\mathbf{G}\mathbf{x} \le \mathbf{h}$
 $\mathbf{A}\mathbf{x} = \mathbf{b}$

$$\min_{\mathbf{x}^+\in\mathbb{R}^n, \mathbf{x}^-\in\mathbb{R}^n, \mathbf{s}\in\mathbb{R}^q} \mathbf{c}^T \mathbf{x}^+ - \mathbf{c}^T \mathbf{x}^- = \tilde{\mathbf{c}}^\top \tilde{\mathbf{x}}$$
s.t. $\mathbf{G}\mathbf{x} + \mathbf{s} = \mathbf{h}$
 $\mathbf{A}\mathbf{x}^+ - \mathbf{A}\mathbf{x}^- = \mathbf{b}$
 $\mathbf{x}^+ \ge 0, \mathbf{x}^- \ge 0, \mathbf{s} \ge 0$

- We express $\mathbf{x} = \mathbf{x}^+ \mathbf{x}^-$ as a difference of positive variables.
- Problem on the right can be reformulated as standard LP.
- The positive variable $s \ge 0$ is used to recover an equality constraint.
- $\blacktriangleright \ \ \text{The standard problem optimizes over } \tilde{\mathbf{x}} = [\mathbf{x}^{+\top}, \mathbf{x}^{-\top}, \mathbf{s}^{\top}]^{\top} \geq \mathbf{0}.$
- \blacktriangleright The matrix $\tilde{\mathbf{A}}$ and $\tilde{\mathbf{b}}$ can be recovered from $\mathbf{A},\,\mathbf{G},\,\mathbf{b}$ and $\mathbf{h}.$
- ▶ The two "tricks" above are classical tools for reformulation.

Primal and Dual problems



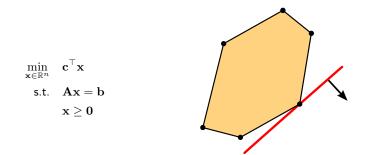
Primal VS Dual

- The problem permute their variables and constraints.
- When there is strict duality (problem has a solution) duality gap is 0:

$$\mathbf{c}^\top \mathbf{x}^\star = -\mathbf{b}^\top \mathbf{v}^\star$$

- Finding x^{*} from v^{*} and vice versa:
 - 1. Find components of \mathbf{x}^{\star} equal to 0 from the equality $(\mathbf{A}^T \mathbf{v}^{\star} \mathbf{c})^T \mathbf{x}^{\star} = 0$.
 - 2. Solve the linear system Ax = b for the non-zero components of x^* .

Solution of the standard LP



Property of the solution

- Problem is convex but possibly has an infinite number of solution (if on one side of the polytope).
- Solution x^{*} is always on a border of the polytop describing the constraints.
- There is at most p ($\mathbf{A} \in \mathbb{R}^{p \times n}$) components of \mathbf{x}^{\star} that are non-zero.
- Those non-zeros components are called active variables.

Robust regression with Least Absolute Deviation

$$\min_{\mathbf{w}\in\mathbb{R}^d} \quad \sum_{i=1}^n |y_i - \mathbf{x}_i^T \mathbf{w}|$$

 More robust to outliers than least squares but also less stable [Barrodale and Roberts, 1973].

Exercise 5: Reformulations as LP

1. Reformulate problem above as a LP with additional variables $e^+ \ge 0, e^- \ge 0$ such that $y - Xw = e^+ - e^-$ with $X = [x_1, \dots, x_n]^T$:

$$\min_{\mathbf{w},\mathbf{e}^+,\mathbf{e}^-}$$

2. Reformulate problem above as a LP with additional variable $f\geq 0_n$ such that $|\mathbf{Hx}-\mathbf{y}|\leq f$:

min w,f

Robust regression with Least Absolute Deviation

$$\min_{\mathbf{w} \in \mathbb{R}^d} \quad \sum_{i=1}^n |y_i - \mathbf{x}_i^T \mathbf{w}|$$

 More robust to outliers than least squares but also less stable [Barrodale and Roberts, 1973].

Exercise 5: Reformulations as LP

1. Reformulate problem above as a LP with additional variables $e^+ \ge 0, e^- \ge 0$ such that $y - Xw = e^+ - e^-$ with $X = [x_1, \dots, x_n]^T$:

$$\min_{\mathbf{w}, \mathbf{e}^+, \mathbf{e}^-} \quad \mathbf{1}_n^T \mathbf{e}^+ + \mathbf{1}_n^T \mathbf{e}^-$$
s.t. $\mathbf{X} \mathbf{w} - \mathbf{y} = \mathbf{e}^+ - \mathbf{e}^-$
 $e^+ \ge \mathbf{0}_n, \mathbf{e}^- \ge \mathbf{0}_n$

2. Reformulate problem above as a LP with additional variable $\mathbf{f} \ge \mathbf{0}_n$ such that $|\mathbf{H}\mathbf{x} - \mathbf{y}| \le \mathbf{f}$: $\min_{\mathbf{w}, \mathbf{f}} \quad \mathbf{1}_n^T \mathbf{f}$ s.t. $\mathbf{X}\mathbf{w} - \mathbf{y} \le \mathbf{f}$, $-\mathbf{X}\mathbf{w} + \mathbf{y} \le \mathbf{f}$ $\mathbf{f} \ge \mathbf{0}_n$

L1 Support Vector Machines

$$\min_{\mathbf{w}\in\mathbb{R}^d} \quad \sum_{i=1}^n \max(0, 1 - y_i \mathbf{x}_i^T \mathbf{w})$$
(12)
s.t. $\|\mathbf{w}\|_1 \le \beta$

Proposed in [Zhu et al., 2004], to promote sparsity in SVM (with the L1 norm).
 Problem above can be reformulated as the following optimization problem :

$$\begin{split} \min_{\mathbf{f}, \mathbf{w}^+, \mathbf{w}^-} \quad \mathbf{1}_n^T \mathbf{f} \\ \text{s.t.} \quad \mathbf{1}_n - (\mathbf{y} \odot \mathbf{X}) \mathbf{w}^+ + (\mathbf{y} \odot \mathbf{X}) \mathbf{w}^- \leq \mathbf{f} \\ \quad \mathbf{1}_d^T \mathbf{w}^+ + \mathbf{1}_d^T \mathbf{w}^- \leq \beta, \quad \mathbf{f} \geq \mathbf{0}, \quad \mathbf{w}^+ \geq \mathbf{0}, \quad \mathbf{w}^- \geq \mathbf{0} \end{split}$$

▶ The corresponding general LP problem with $\mathbf{x} = [\mathbf{w}^{+T}, \mathbf{w}^{-T}, \mathbf{f}]^T$ has the following matrices:

$$\mathbf{c} = \begin{bmatrix} \mathbf{0} \\ \mathbf{0} \\ \mathbf{1}_n \end{bmatrix}, \quad \mathbf{G} = \begin{bmatrix} -(\mathbf{y} \odot \mathbf{X}) & (\mathbf{y} \odot \mathbf{X}) & -\mathbf{I}_n \\ \mathbf{1}_{1,d} & \mathbf{1}_{1,d} & \mathbf{0}_{1,n} \\ -\mathbf{I}_d & \mathbf{0}_{d,d} & \mathbf{0}_{d,n} \\ \mathbf{0}_{d,d} & -\mathbf{I}_d & \mathbf{0}_{d,n} \\ \mathbf{0}_{n,d} & \mathbf{0}_{n,d} & -\mathbf{I}_n \end{bmatrix}, \qquad \mathbf{h} = \begin{bmatrix} -\mathbf{1}_n \\ \beta \\ \mathbf{0}_d \\ \mathbf{0}_d \\ \mathbf{0}_n \end{bmatrix}$$

6.2.2 - Linear Program (LP) - Examples of LP in ML - 27/52

Simplex Algorithm

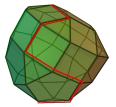
Main idea behind the simplex

- Initialize with a basic feasible solution x⁽⁰⁾ (on a vertex or extreme point of the polytope).
- Update the solution to decrease the loss at each iteration.
- Use the sparsity of x (add and remove active variables).

Simplex algorithm

- Invented by Dantzig around 1957.
- Solved the problem he thought was a homework exercise from his course.
- Standard algorithm for solving LP, very efficient for sparse problems but possibly non polynomial (worst case).
- On network flow problems, the adapted network simplex is proven to be polynomial [Orlin, 1997] (optimal transport).
- in Python : scipy.optimize.linprog(method='simplex')

More details: [Vanderbei et al., 2015, part 1]



Interior point solver

Interior point method (IPM) for LP

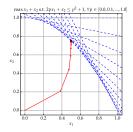
$$\min_{\mathbf{x} \in \mathbb{R}^n} \mathbf{c}^\top \mathbf{x}$$
s.t. $\mathbf{G}\mathbf{x} < \mathbf{h}$

- Simplex searches on the corners of the polytope, IPM optimize inside it.
- Never against the constraints until numerical precision is achieved.
- Polynomial complexity for LP (better than simplex in theory).
- In Python: scipy.optimize.linprog

More details: [Boyd and Vandenberghe, 2004, Chapter 11], [Vanderbei et al., 2015, Part 3], [Nocedal and Wright, 2006, Chapter 14]

 \Rightarrow

$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad \delta \mathbf{c}^\top \mathbf{x} + -\sum_{i=1}^q \log(\mathbf{g}_i^T \mathbf{x} - h_i)$$



Solving a Linear Program

Simplex and variants

- Exact solutions.
- Can be slow for large problems.
- Use it on structured graph flow.

LP solvers in Python

Interior point problem

- Better at early stopping.
- Usually faster on large problems.
- Most generic commercial solvers.
- Scipy : scipy.optimize.linprog function (both simplex and interior points)
- cvxopt : Interior point solver for standard problems + wrapper for commercial solvers and GLPK [Vandenberghe, 2010].
- Mosek Commercial solver (free for academics) [Andersen and Andersen, 2000].
- **Gurobi** Commercial solver (free for academics).
- CPLEX Commercial solver (free for academics).

Wrappers available : https://github.com/stephane-caron/lpsolvers

Quadratic Program

Optimization problem

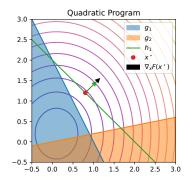
$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad \frac{1}{2}\mathbf{x}^T\mathbf{Q}\mathbf{x} + \mathbf{c}^T\mathbf{x} \qquad (13)$$

s.t. $\mathbf{G}\mathbf{x} \leq \mathbf{h}$
 $\mathbf{A}\mathbf{x} = \mathbf{b}$

- ▶ Q ∈ ℝn × n is a symmetric positive definite matrix (convex QP).
- $\mathbf{A} \in \mathbb{R}^{p \times n}, \mathbf{b} \in \mathbb{R}^{p}$, and $\mathbf{G} \in \mathbb{R}^{q \times n}, \mathbf{h} \in \mathbb{R}^{q}$.
- Most standard solvers (open source and commercial) use this formulation.

Special cases

- Unconstrained : close form solution or iterative methods (Conjugate gradients)
- Box constraints $l \le x \le u$: projected gradient (see proximal methods).



QP Exemple: Portfolio optimization

- Model proposed by Markowitz in 1952 (Nobel Prize in economy).
- x is a portfolio of n assets (or stock).
- The price change for each asset is modeled as random variables with expected price change p and covariance Σ.
- For a given portfolio x
 - The expected gain (return) is : $\mathbf{p}^T \mathbf{x}$
 - The expected variance is : $\mathbf{x}^T \mathbf{\Sigma} \mathbf{x}$
- The portfolio optimization can be expressed for a positive balance b > 0 as:

$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad \mathbf{x}^T \mathbf{\Sigma} \mathbf{x} \tag{14}$$

s.t.
$$\mathbf{1}_n^T \mathbf{x} = b$$
 (15)

$$\mathbf{p}^T \mathbf{x} \ge r_{min} \tag{16}$$

where r_{min} is the minimal return of the portfolio.

Special Case : QP without constaints

$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad \frac{1}{2}\mathbf{x}^T\mathbf{Q}\mathbf{x} + \mathbf{c}^T\mathbf{x}$$
(17)

Unconstrained QP

- ▶ The gradient of the term above is $\nabla_{\mathbf{x}} = \frac{1}{2} (\mathbf{Q} + \mathbf{Q}^T) \mathbf{x} + \mathbf{c}$
- \blacktriangleright For symmetric matrix ${\bf Q}$ a solution respects : ${\bf Q}{\bf x}^{\star}=-{\bf c}$
- $\blacktriangleright\,$ If ${\bf Q}$ is invertible and strictly positive definite then : ${\bf x}^{\star}=-{\bf Q}^{-1}{\bf c}$
- To solve the problem several approaches
 - 1. Solve the linear equations : np. linalg . solve with complexity $O(n^3)$
 - 2. Solve the linear equations with Conjugate Gradient or other gradient descent methods (see other courses).

Exercise 6: Least Square and Ridge

$$\min_{\mathbf{x}} \quad \frac{1}{2} \|\mathbf{H}\mathbf{x} - \mathbf{y}\|^2 \qquad \qquad \min_{\mathbf{x} \in \mathbb{R}^n} \quad \frac{1}{2} \|\mathbf{H}\mathbf{x} - \mathbf{y}\|^2 + \lambda \frac{1}{2} \|\mathbf{x}\|^2$$

Recover the matrices ${\bf Q}$ and ${\bf c}$ of the equivalent QP for the problems above:

$$\mathbf{Q} = \mathbf{c} = \mathbf{Q} = \mathbf{c} =$$

Special Case : QP without constaints

$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad \frac{1}{2}\mathbf{x}^T\mathbf{Q}\mathbf{x} + \mathbf{c}^T\mathbf{x}$$
(17)

Unconstrained QP

- ▶ The gradient of the term above is $\nabla_{\mathbf{x}} = \frac{1}{2}(\mathbf{Q} + \mathbf{Q}^T)\mathbf{x} + \mathbf{c}$
- For symmetric matrix \mathbf{Q} a solution respects : $\mathbf{Q}\mathbf{x}^{\star} = -\mathbf{c}$
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- To solve the problem several approaches
 - 1. Solve the linear equations : np. linalg . solve with complexity $O(n^3)$
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Exercise 6: Least Square and Ridge

$$\min_{\mathbf{x}} \quad \frac{1}{2} \|\mathbf{H}\mathbf{x} - \mathbf{y}\|^2 \qquad \qquad \min_{\mathbf{x} \in \mathbb{R}^n} \quad \frac{1}{2} \|\mathbf{H}\mathbf{x} - \mathbf{y}\|^2 + \lambda \frac{1}{2} \|\mathbf{x}\|^2$$

Recover the matrices ${\bf Q}$ and ${\bf c}$ of the equivalent QP for the problems above:

$$\mathbf{Q} = \mathbf{H}^T \mathbf{H}, \qquad \mathbf{c} = -\mathbf{H}^T \mathbf{y}, \qquad \qquad \mathbf{Q} = \mathbf{H}^T \mathbf{H} + \lambda \mathbf{I}, \qquad \mathbf{c} = -\mathbf{H}^T \mathbf{y},$$

6.3.1 - Quadratic Program (QP) - Problem formulation - 33/52

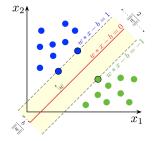
Support Vector Machines (1)

Hard margin SVM [Cortes and Vapnik, 1995]

$$\min_{\mathbf{w},b} \quad \frac{1}{2} \|\mathbf{w}\|^2$$
s.t. $y_i(\mathbf{x}_i^T \mathbf{w} + b) \ge 1$

$$(18)$$

- All samples (x_i, y_i) must be classified well with a margin of at least 1.
- Needs the data to be linearly separable.
- Minimizing the norm of w corresponds to maximizing the margin ²/_{||w||}.



Soft margin SVM

$$\min_{\mathbf{w}\in\mathbb{R}^{d},b\in\mathbb{R}} \quad C\sum_{i} \max(0,1-y_{i}(\mathbf{x}_{i}^{T}\mathbf{w}+b)) + \frac{1}{2}\|\mathbf{w}\|^{2}$$
(19)

- The margin constraints are relaxed with the Hinge loss.
- C is the weight of the data fitting term.
- Non differentiable convex problem.

Support Vector Machines (2)

Primal SVM

$$\min_{\mathbf{w}\in\mathbb{R}^{d}, b\in\mathbb{R}, \mathbf{z}\in\mathbb{R}^{n}} \quad C\sum_{i} z_{i} + \frac{1}{2} \|\mathbf{w}\|^{2}$$
s.t. $y_{i}(\mathbf{x}_{i}^{T}\mathbf{w} + b) \geq 1 - z_{i}, \forall i$
 $\mathbf{z} \geq \mathbf{0}$

$$(20)$$

• We introduce the variables $z_i \ge 0$ such that $z_i = \max(0, 1 - y_i(\mathbf{x}_i^T \mathbf{w} + b))$.

Dual SVM

$$\min_{\boldsymbol{\alpha} \in \mathbb{R}^{n}} \quad \frac{1}{2} \boldsymbol{\alpha}^{T} \mathbf{Q} \boldsymbol{\alpha} - \mathbf{1}_{n}^{T} \boldsymbol{\alpha}$$
(21)
s.t. $\mathbf{y}^{T} \boldsymbol{\alpha} = 0$
 $\mathbf{0}_{n} \leq \boldsymbol{\alpha} \leq C \mathbf{1}_{n}$

▶ QP $(Q_{i,j} = y_i y_j \mathbf{x}_i^T \mathbf{x}_j)$ with box constraints and one linear constraint.

- Primal solution can be recovered with : $\mathbf{w}^* = \sum_i y_i \alpha_i^* \mathbf{x}_i$.
- b^{*} can be found on a support vector where inequality becomes equality.
- Most common formulation because allows the use of kernel for nonlinear classification (Q_{i,j} = y_iy_jk(x_i, x_j))

Lasso estimator

$$\min_{\mathbf{w}} \quad \frac{1}{2} \|\mathbf{X}\mathbf{w} - \mathbf{y}\|^2 + \lambda \sum_i |w_i|$$
(22)

- Classical approach to perform regression with variable selection [Tibshirani, 1996].
- Quadratic data fitting, L1 regularization term.
- Expressed either as additive term or constraint (equivalent problem).

Exercise 7: Lasso reformulation as QP

1. Reformulate the Lasso problem as a positive QP with $\mathbf{w}=\mathbf{w}^+-\mathbf{w}^-$

$$\begin{array}{ll} \min_{\mathbf{w}^+,\mathbf{w}^-} \\ \text{s.t.} \quad \mathbf{w}^+ \geq \mathbf{0}, \ \mathbf{w}^- \geq \mathbf{0} \end{array}$$

2. Express the matrices Q, c, G, h for standard QP corresponding to the problem.

$$\mathbf{Q} = \mathbf{c} = \mathbf{G} = \mathbf{h} =$$

Lasso estimator

$$\min_{\mathbf{w}} \quad \frac{1}{2} \|\mathbf{X}\mathbf{w} - \mathbf{y}\|^2 + \lambda \sum_i |w_i|$$
(22)

Classical approach to perform regression with variable selection [Tibshirani, 1996].

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$$\begin{split} \min_{\mathbf{w}^+,\mathbf{w}^-} & \frac{1}{2} \|\mathbf{X}(\mathbf{w}^+ - \mathbf{w}^-) - \mathbf{y}\|^2 + \lambda \sum_i w_i^+ + w_i^- \\ \text{s.t.} & \mathbf{w}^+ \geq \mathbf{0}, \ \mathbf{w}^- \geq \mathbf{0} \end{split}$$

2. Express the matrices Q, c, G, h for standard QP corresponding to the problem.

$$\mathbf{Q} = \mathbf{c} = \mathbf{G} = \mathbf{h} =$$

Lasso estimator

1

$$\min_{\mathbf{w}} \quad \frac{1}{2} \|\mathbf{X}\mathbf{w} - \mathbf{y}\|^2 + \lambda \sum_i |w_i|$$
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2. Express the matrices $\mathbf{Q}, \mathbf{c}, \mathbf{G}, \mathbf{h}$ for standard QP corresponding to the problem.

$$\mathbf{Q} = \begin{bmatrix} \mathbf{X}^T \mathbf{X} & -\mathbf{X}^T \mathbf{X} \\ -\mathbf{X}^T \mathbf{X} & \mathbf{X}^T \mathbf{X} \end{bmatrix}, \qquad \mathbf{c} = \begin{bmatrix} -\mathbf{X}^T \mathbf{y} + \lambda \mathbf{1}_d \\ \mathbf{X}^T \mathbf{y} + \lambda \mathbf{1}_d \end{bmatrix}, \qquad \mathbf{G} = -\mathbf{I}_{2d}, \qquad \mathbf{h} = \mathbf{0}_{2d}$$

Active set Algorithm

$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad \frac{1}{2}\mathbf{x}^T\mathbf{Q}\mathbf{x} + \mathbf{c}^T\mathbf{x} \\ s.t. \quad \mathbf{G}\mathbf{x} \le \mathbf{h} \\ \mathbf{A}\mathbf{x} = \mathbf{b}$$

Principle of active set method

- Search for the active constraints $\mathcal{A}(\mathbf{x}^*)$.
- If the optimal active set is known the problem is an equality constrained QP.
- ▶ QP with equality constraint can be solved with : null space + unconstrained QP.
- QP version of the simplex (search on which constraints is the solution).
- Very efficient on some problems (positivity, bloc constraints, SVM).

Active set Method (simplified)

Initialize feasible x , $\mathcal{A}(\mathbf{x}) = \{i | \mathbf{g}_i^T \mathbf{x} = h_i\}$ the active set of inequality constraints.

- 1. Solve subproblem with inequality constraints in $\mathcal{A}(\mathbf{x})$ forced to equality.
- 2. Update the active set using KKT conditions.

More details: [Nocedal and Wright, 2006, Sec. 16.5]

Sequential Minimal Optimization (SMO)

$$\min_{\boldsymbol{\alpha} \in \mathbb{R}^n} \quad \frac{1}{2} \boldsymbol{\alpha}^T \mathbf{Q} \boldsymbol{\alpha} - \mathbf{1}_n^T \boldsymbol{\alpha}$$

s.t. $\mathbf{y}^T \boldsymbol{\alpha} = 0$
 $\mathbf{0}_n \le \boldsymbol{\alpha} \le C \mathbf{1}_n$

Principle of SMO

- Proposed by [Platt, 1998] to solve large scale SVM.
- Coordinate descent algorithm taking into account $\mathbf{y}^T \boldsymbol{\alpha} = 0$.
- The choice of the coordinates to update is sensitive.
- Sate of the art solver for SVM [Chang and Lin, 2001] also use a cache for computing the kernel matrix.

SMO Algorithm

Initialize feasible lpha

- 1. Find two components α_i and α_j that violate KKT conditions.
- 2. Solve the QP on only those components in closed form (1D problem).

Solving a QP

$$\min_{\mathbf{x} \in \mathbb{R}^n} \quad \frac{1}{2} \mathbf{x}^T \mathbf{Q} \mathbf{x} + \mathbf{c}^T \mathbf{x}$$

s.t. $\mathbf{G} \mathbf{x} \le \mathbf{h}$
 $\mathbf{A} \mathbf{x} = \mathbf{b}$

Main Algorithms

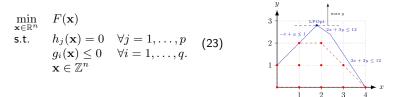
- Interior points Efficient for large problems (commercial solvers).
- Active set General solver, an be very fast on structured problems (sparsity, SVM)
- SMO State of the art solver for SVM.

QP Solvers in Python

- Numpy (no constraints): (np. linalg . solve ornp. linalg . lstsq).
- quadprog : Implements active set [Goldfarb and Idnani, 1983]
- **cvxopt** : Interior point solver for standard problems + wrapper for Mosek.
- OSQP : Operator spliting QP solver [Stellato et al., 2017].
- Mosek : Commercial solver (free for academics) [Andersen and Andersen, 2000].
- **Gurobi** : Commercial solver (free for academics).

Benchmark available : https://github.com/qpsolvers/qpbenchmark 6.3.3 - Quadratic Program (QP) - Active set (AS) and Sequentiel Minimal Optimization (SMO) - 39/52

Integer Programming



• Classical optimization problem with additional integer constraints $\mathbf{x} \in \mathbb{Z}^n$.

- Zero-one programming when variables can be only binary $\mathbf{x} \in \{0,1\}^n$.
- ► Mixed Integer Programming (MIP) problems when only part of the variables are integer : x_i ∈ Z for i = 1,..., n_i with n_i ≤ n.
- Problem is extremely hard to solve exactly (NP complete).

Algorithms

- Continuous relaxation (and then rounding, can work well on MILP).
- Cutting Plane Algorithm (relaxation + iteratively add linear constraints).
- Branch and bound (exact method using upper and lower bounds to split the space of solution).

MILP and MIQP

Mixed Integer LP (MILP)Mixed Integer QP (MIQP) $\min_{\mathbf{x} \in \mathbb{R}^n}$ $\mathbf{c}^T \mathbf{x}$ $\min_{\mathbf{x} \in \mathbb{R}^n}$ $\frac{1}{2} \mathbf{x}^T \mathbf{Q} \mathbf{x} + \mathbf{c}^T \mathbf{x}$ s.t. $\mathbf{x} \ge \mathbf{0}$ s.t. $\mathbf{x} \ge \mathbf{0}$ $\mathbf{A}\mathbf{x} = \mathbf{b}$ $\mathbf{A}\mathbf{x} = \mathbf{b}$ $\mathbf{A}\mathbf{x} = \mathbf{b}$ $x_i \in \mathbb{Z}, \forall i \in \{1, \dots, n_i\}$ $x_i \in \mathbb{Z}, \forall i \in \{1, \dots, n_i\}$

- Well studied MIP problems.
- For MILP, relaxation can be exact (when total unimodularity of A)
- Solved by Branch and Bound and cutting planes in general.

MIP solvers in Python

- **cvxpy** : General optimization (multiple wrappers) [Diamond and Boyd, 2016].
- **ECOS** : Embedded Conic Solver for MILP [Domahidi et al., 2013].
- Mosek : Commercial solver (free for academics) [Andersen and Andersen, 2000].
- Gurobi : Commercial solver (free for academics).

L0 sparse regression

$$\min_{\mathbf{x} \in \mathbb{R}^n} \quad \frac{1}{2} \|\mathbf{H}\mathbf{x} - \mathbf{y}\|^2 + \lambda \|\mathbf{x}\|_0$$

Problem above can be reformulated as a MIQP [Bourguignon et al., 2015].

First we introduce a binary vector $\mathbf{z} \in \{0, 1\}^n$.

2

• We suppose that $z_i = 1$ if variable $\mathbf{x}_i \neq 0$ else $z_i = 0$. This means that for a big enough M we have:

$$-M\mathbf{z} \le \mathbf{x} \le M\mathbf{z}$$

We can express the L0 sparse regression as the following optimization problem:

$$\min_{\mathbf{x}\in\mathbb{R}^{n},\mathbf{z}\in\mathbb{R}^{n}} \quad \frac{1}{2}\mathbf{x}^{\top}\mathbf{H}^{T}\mathbf{H}\mathbf{x} - (\mathbf{H}^{T}\mathbf{y})^{T}\mathbf{x} + \lambda \mathbf{1}_{n}^{T}\mathbf{z}$$

s.t. $-M\mathbf{z} \leq \mathbf{x} \leq M\mathbf{z}$
 $\mathbf{z} \in \{0,1\}^{n}$

Other formulations corresponds to constrained expression but all use the "big M" trick.

Quadratically Constrained QP (QCQP)

Optimization problem

$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad \frac{1}{2}\mathbf{x}^T\mathbf{Q}_0\mathbf{x} + \mathbf{c}_0^T\mathbf{x} \tag{24}$$

s.t.
$$\mathbf{x}^T \mathbf{Q}_i \mathbf{x} + \mathbf{c}_i^T \mathbf{x} \le \mathbf{h}_i, \ \forall i = 1, \dots, m$$

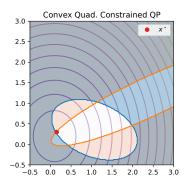
 $\mathbf{A}\mathbf{x} = \mathbf{b}$

- If Q₀,..., Q_m are positive definite then the problem is convex and can be solved with interior point.
- ► Nonconvex QCQP is NP-hard, because a constraint x_i ∈ {0,1} is recovered with:

$$x_i(1-x_i) \ge 0$$
 and $x_i(1-x_i) \le 0$

QCQP solvers in Python

- cvxpy : with nonconvex QCQP extension [Park and Boyd, 2017] .
- Mosek : Commercial solver (free for academics) [Andersen and Andersen, 2000].
- Gurobi : Commercial solver (free for academics).



Quadratically Constrained QP (QCQP)

Optimization problem

$$\min_{\mathbf{x}\in\mathbb{R}^n} \quad \frac{1}{2}\mathbf{x}^T\mathbf{Q}_0\mathbf{x} + \mathbf{c}_0^T\mathbf{x} \tag{24}$$

s.t.
$$\mathbf{x}^T \mathbf{Q}_i \mathbf{x} + \mathbf{c}_i^T \mathbf{x} \le \mathbf{h}_i, \ \forall i = 1, \dots, m$$

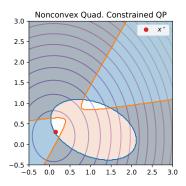
 $\mathbf{A}\mathbf{x} = \mathbf{b}$

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K-means as MIQCQP

$$\min_{\bar{\mathbf{x}}_k, \forall_k} \quad \sum_{i=1}^N \min_k \|\bar{\mathbf{x}}_k - \mathbf{x}_i\|^2$$

- ▶ The argmin for each sample can be replaced by a linear term with a matrix $\mathbf{Z} \in \{0, 1\}^{N, K}$ modeling the clustering of the samples.
- We force a unique cluster selection with constraints

$$\mathbf{Z} \in \{0,1\}^{N,K}, \quad \mathbf{Z}\mathbf{1}_K = \mathbf{1}_N$$

- We introduce the distance variable as $D_{i,k} = \|\mathbf{x}_i \bar{\mathbf{x}}_k\|^2$
- The optimization problem above can be expressed as

$$\min_{\bar{\mathbf{x}}_{k}, \forall_{k}, \mathbf{Z} \in \mathbb{R}^{N \times K}, \mathbf{D} \in \mathbb{R}^{N \times K}} \sum_{i,k} Z_{i,k} D_{i,k} \qquad (25)$$
s.t. $D_{i,k} = \|\mathbf{x}_{i} - \bar{\mathbf{x}}_{k}\|^{2}, \forall_{i}, \forall k$
 $\mathbf{Z} \mathbf{1}_{K} = \mathbf{1}_{N}$
 $\mathbf{Z} \in \{0, 1\}^{N, K}$

Warning: Never try to solve K-means with this formulation!

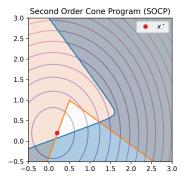
Second Order Cone Programming (SOCP)

Optimization problem

 $\min_{\mathbf{x}\in\mathbb{R}^n} \mathbf{c}^T \mathbf{x}$ (26)

- s.t. $\|\mathbf{A}_i \mathbf{x} \mathbf{b}_i\|_2 \leq \mathbf{h}_i^T \mathbf{x} + d_i, \quad i = 1, \dots, m$ $\mathbf{A}_0 \mathbf{x} = \mathbf{b}_0$
- The following constraint is called a Second order cone constraint:

$$\|\mathbf{A}\mathbf{x} - \mathbf{b}\|_2 \le \mathbf{h}^T \mathbf{x} + d$$



- When $\mathbf{h}_i = \mathbf{0}, \ \forall i$ the problem is a QCQP (one can square the norm).
- Other kind of cone constraints can be used (positive definite matrices).

SOCP solvers in Python

- **cvxopt** : Interior point solver [Vandenberghe, 2010].
- cvxpy : SOCP solver [Diamond and Boyd, 2016].
- Mosek : Commercial solver (free for academics) [Andersen and Andersen, 2000].
- **Gurobi** : Commercial solver (free for academics).

Robust Support Vector Machines

$$\min_{\mathbf{w}\in\mathbb{R}^{d},b\in\mathbb{R},\mathbf{z}\in\mathbb{R}^{n}} C\sum_{i} z_{i} + \frac{1}{2} \|\mathbf{w}\|^{2}$$
s.t. $y_{i}(\mathbf{x}_{i}^{T}\mathbf{w} + b) \geq 1 - z_{i} + \gamma_{i} \left\|\boldsymbol{\Sigma}_{i}^{\frac{1}{2}}\mathbf{w}\right\|, \forall i$
 $\mathbf{z} \geq \mathbf{0}$

$$(27)$$

- Proposed in [Shivaswamy et al., 2006] to handle uncertain and missing data.
- We suppose that we have uncertain data (x_i, y_i) and that the training sample x_i has a covariance matrix Σ_i to model its uncertainty.
- In this can one want to replace the hard margin constraint by a probabilistic variant

$$P(y_i(\mathbf{x}_i^T\mathbf{w}+b) \ge 1-z_i) \ge 1-\kappa_i$$

were κ_i is small.

• When using the normal distribution on the training samples, one can recover the optimization problem above with $\gamma_i = \phi^{-1}(\kappa_i)$ where ϕ is the normal CDF.

Semi-Definite Programming

Optimization problem

$$\begin{split} \min_{\mathbf{X} \in \mathbb{S}^n} & \langle \mathbf{X}, \mathbf{C} \rangle_{\mathbb{S}^n} \\ \text{s.t.} & \langle \mathbf{X}, \mathbf{A}_i \rangle_{\mathbb{S}^n} = b_i, \quad i = 1, \dots, m \\ & \mathbf{X} \succeq 0 \end{split}$$

- \mathbb{S}^n is the set of $n \times n$ symmetric matrices.
- ▶ $\langle \mathbf{X}, \mathbf{C} \rangle_{\mathbb{S}^n} = \sum_{i,j} X_{i,j} C_{i,j}$ is the Frobenius scalar product between matrices.
- The constraint $\mathbf{X} \succeq 0$ forces \mathbf{X} to be semi-definite positive.
- Special case of cone programming (cone of positive semi-definite matrices).
- Can be solved efficiently with interior point solver.

SDP solvers in Python

- cvxopt : Interior point solver [Vandenberghe, 2010].
- cvxpy : SDP solver [Diamond and Boyd, 2016].
- Mosek : Commercial solver (free for academics) [Andersen and Andersen, 2000].
- **Gurobi** : Commercial solver (free for academics).

(28)

Conclusion

Constrained optimization

- Constrained optimization is a very large field.
- Langrangian and KKT conditions are the main tools to solve and check solutions.

Standard Problems (properties)

- Important properties
 - Linear or quadratic objective function.
 - Linear, quadratic of conic constraints.
 - Real of integer variables.
- Many existing generic solvers for those problems (commercial or free).

Solving DS and ML problems beyond gradient descent

- Constraints are important in many problems (fairness, robustness).
- ▶ You often have to model new ML problems depending on the constraints.
- Reformulation is key to use generic solvers (important skill).

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